

QUIC: the details

Robin Marx - **@programmingart** PhD researcher – Hasselt University



https://quic.edm.uhasselt.be Curl-up – Prague – March 2019

QUIC in Theory and Practice @ DeltaV 2018





https://www.youtube.com/watch?v=B1SQFjIXJtc http://internetonmars.org/deltav.pdf My self-worth is determined by random, white men on the internet

∥ 3



Richard M. 9 months ago

Using comic book characters and attempts to be funny fell flat. Don't do such things when your audience may know little about comics and are looking for technical info. You lost me one or two minutes in.

1 2 👎 REPLY

James Fuller

to me 👻

The curl up day (Sat/Sun) are 'turn it up to 11' in terms of technical level,



QUIC is "quite" complex...

- 6 "Core" specifications:
 - QUIC invariants : 9 pages
 - QPACK header compression : 37 pages
 - Recovery: loss detection and congestion control : 38 pages
 - TLS mapping : 43 pages
 - HTTP/3 : 53 pages
 - Core transport : 139 pages

319 pages total

(9 more than The Hobbit)

2566 issues on github

- At least 20 more side-specs
 - Applicability, manageability, spinbit, DATAGRAM, ...



https://github.com/quicwg/base-drafts

1. Connection setup: Similar to TLS 1.3, but small differences Client Server ClientHello (0-RTT Application Data) ----> ServerHello {EncryptedExtensions} {Finished} No "End-of-early-data" [Application Data] No Record layer {Finished} ----> [Application Data] [Application Data] <----> 0-RTT is done in QUIC itself QUIC performs all actual () Indicates messages protected by early data (0-RTT) keys encrypt/decrypt operations {} Indicates messages protected using handshake keys Indicates messages protected using application data Γ1 (1-RTT) keys

Figure 1: TLS Handshake with 0-RTT

1. Connection setup: TLS is (largely) abstracted out

►► UHASSELT EDM



Figure 19: CRYPTO Frame Format

TLS is just another opaque bytestream

cryptoData = TLS.getTLSstuff()
sendCryptoFrame(cryptoData)

TLS.processTLSstuff(cryptoData) TLS.onNewKeyAvailable(keyCallback) 1. Connection setup: TLS is (largely) abstracted out: CRYPTO frames

RX Initial(0x0), Version: 0xff00000f, Dest CID: 0x1156a3c691c16e8c390499d90aac254b21, Src CID: 0x0736553ca2daa8 PKN: 0, payload length: 1184 CRYPTO (0x18) length=304 offset=0 data=SSL3_MT_CLIENT_HELLO (PROBABLY, but could be different or more if split or coalesced packet) PADDING (0×0) length= 859 , Version: 0xff00000f, Dest CID: 0x0736553ca2daa8, Src CID: 0x0f72f17fbaab34b215b0990e570661 , payload length: 1220 length=1200 offset=0 data=SSL3_MT_ENCRYPTED_EXTENSIONS (PROBABLY, but could be different or more if split or coalesced packet)

►► UHASSELT EDM

https://github.com/rmarx/quicker



Table 1: Encryption Levels by Packet Type





▶ UHASSELT EDM

https://github.com/quicwg/base-drafts/issues/1018

1. Connection setup: Each encryption level has a separate packet type too

In this version of QUIC, the following packet types with the long header are defined:

_		L
Type	Name	Section
0x0	Initial	Section 17.2.2
0×1	0-RTT	<u>Section 17.2.3</u>
0x2	Handshake	<u>Section 17.2.4</u>
0x3	Retry	<u>Section 17.2.5</u>
T		

Table 5: Long Header Packet Types

Only 1-RTT packets use the "short" header



1. Long packet headers: lots of information

```
-+-+-+-+-+-+
111 2 | R R | P P
              -+-+-+-+-+-+-+-+-+-+-+
Version (32)
          -+-+-+-+-+-+-+-+-+-+
DCIL(4) SCIL(4)
        +-+-+-+-+-+-
       Destination Connection ID (0/32..144)
    Source Connection ID (0/32..144)
                              . . .
       Length (i)
     Packet Number (8/16/24/32)
                              . . .
          -+-+-+-+-+-+-+-+-+-+
             Payload (*)
```

Figure 12: <u>Handshake</u> Protected Packet

1. Long packet headers: information that never changes afterwards



Figure 12: <u>Handshake</u> Protected Packet

►► UHASSELT EDM

1. Short packet headers: optimized headers for most data packets

17.3. Short Header Packets

This version of QUIC defines a single packet type which uses the short packet header.



Figure 14: Short Header Packet Format

►► UHASSELT EDM

1. Header protection : actually more than 4 keys!

N different "keys"

- Initial (AEAD) : Can't be changed, but can be read
- O-RTT (pre-shared) : Can't be changed or read, but replayed
- Handshake
- 1-RTT
- Header Protection key (same for all packet types)
 - Protects flags + Packet Number (NOT connectionID!)
- Unlimited amount of new 1-RTT keys
 - KEY_PHASE bit in the flags



2. ConnectionID



- Why do we need it?
- Why do we need 2? And then back to 1?
- Why is it so large?

►► UHASSELT EDM

2. ConnectionID: Connection migration and NAT rebinding



 4G <-> Wifi
 Network Address Translation

Both cases can change IP and/or port

ConnectionID can stay the same (or change to previously agreed upon values) 2. ConnectionID: Asymmetric for Routing/Load balancing

Originally: just 1 ConnectionID Later: source vs destination

Want server to choose CID, to use for routing/load balancing

- Need traffic to go to same origin server
- Put routing info inside the CID for stateless load balancers



Statefull:

Remember that 0xabcdefab goes to server 1

Stateless:

Oxabcdefa2 always goes to server 2

- \rightarrow Client doesn't know this
- \rightarrow Server 2 has to choose its own CID

https://datatracker.ietf.org/doc/draft-duke-quic-load-balancers/ image: http://tutorials.jenkov.com/software-architecture/load-balancing.html

2. ConnectionID: Client chooses for server, server overrides

RX Initial(0 PKN: 0, CRYPTO (0x length=304 data=SSL3_ PADDING (0 length= 85	<pre>x0), Version: 0xff0 payload length: 11 18) offset=0 MT_CLIENT_HELLO (PR x0) 9</pre>	0000f, Dest CID: 84 OBABLY, but could	0x1156a3c691c1 d be different	6e8c390499d90aad	254b21, Src CI t or coalesced	D: 0x0736553ca2daa8 packet)
TX Handshake PKN: 0 CRYPTO (0) length=120 data=SSL3	(0x2), Version: 0x1 , payload length: 12 (18) 00 offset=0 _MT_ENCRYPTED_EXTENS	ff00000f, Dest C 220 SIONS (PROBABLY,	ID: 0x0736553ca but could be c	2daa8, Src CID:	0x0f72f17fbaal e if split or o	b34b215b0990e570661 coalced packet)
►► UHASSE	LT EDM					https://github.com/rmarx/quicker

2. ConnectionID: Generic metadata storage field

Load balancer wants to support changing CIDs

 \rightarrow Same routing info should be encoded in different CIDs

- \rightarrow Cannot do this for just random values
- \rightarrow CID is routing metadata but encrypted!
- \rightarrow Metadata can be large, so CID up to 18 bytes!

Facebook also encodes process ID

UHASSELT EDM

 \rightarrow Seamless handover on server upgrade



First CID:

Decrypted: 0xabcdefa2 Encrypted: 0x12345678

After CID change: Decrypted: 0xaabbccd2 Encrypted: 0x87654321

https://datatracker.ietf.org/doc/draft-duke-quic-load-balancers/ https://conferences2.sigcomm.org/co-next/2018/slides/epiq-keynote.pdf https://www.youtube.com/watch?v=8IYHNzoPS20

3. Minimize overhead

Transmission Control Protocol (TCP) Header 20-60 bytes



- Fields are always the same size (even if values are small)
- Good for fast processing, bad for minimizing overhead
- 1st example was Short header, there are many more!

►► UHASSELT EDM

3. Minimize overhead: QUIC saves on bits like crazy!

Variable Length Integer Encoding + optional fields





- 3. Minimize overhead: But not always in the same way...
- Not everything is VLIE though



- DCIL : Destination Connection ID Length
 - Faster parsing (e.g., routing, load balancing)
- DCIL = 0 : no DCID set
- For all other values: have to do +3!

0000 = 0 bytes 0001 = 4 bytes 0010 = 5 bytes

1111 = 18 bytes



3. Minimize overhead: It gets worse

- Packet numbers are ~delta-encoded
 - Even large packet numbers use only a small amount of bytes
 - More packets on the wire = more bytes needed
 - Typically 1 or 2 is enough!

UHASSELT | EDN

- Done by "delta-encoding" based on the largest acknowledged packet
- Complex algorithm that allows leaving out several least-significant bits as well



- Packet number length in 2 flag bits (used to be VLIE though)
 - 2 bits, number always at least 1 byte long, so always +1:
 - 00 = 1 byte, 01 = 2 bytes, 10 = 3 bytes, 11 = 4 bytes

3. Minimize overhead: yet another trick up my sleeve



Figure 17: ACK Block Section

UHASSELT EDM

Full packet number

Not full numbers, but COUNTS, relative to the previous values - Counts stay very small

"The number of packets in the gap is one higher than the encoded value of the Gap Field"

Most of these will fit in 1 byte each

https://tools.ietf.org/html/draft-ietf-quic-transport-19#appendix-A

3. Minimize overhead: yet another trick up my sleeve: Reprise

- ACK delay: measured in microseconds
 - Indicates the processing delay to generate the ACK
 - Allows better estimates for network delay (total processing = network)
 - Maybe not everyone needs this high resolution...
 - Value in the packet needs to be multiplied by 2[^]ack_delay_exponent
 - Peer can indicate ack_delay_exponent in transport parameters
 - Default: 2^3
 - E.g., delay was 400 microseconds, value on the wire is 50
 - Anything lower than 8 microseconds cannot be encoded



https://tools.ietf.org/html/draft-ietf-quic-transport-19#appendix-A

4. HTTP/3 vs HTTP/2 : Some simplification

- QUIC takes over several things that were in H2
 - Flow Control

UHASSELT EDM

- Streams and stream management
- Padding frames
- Fun fact: without HTTP/3 implementations, QUIC was tested with HTTP/0.9

> https://tools.ietf.org/html/draft-ietf-quic-transport-19 https://tools.ietf.org/html/draft-ietf-quic-http-19 https://github.com/rmarx/quicker



https://www.youtube.com/watch?v=B1SQFjIXJtc http://internetonmars.org/deltav.pdf

▶ UHASSELT EDM



http://internetonmars.org/deltav.pdf

►► UHASSELT EDM



4. HTTP/3 vs HTTP/2 : Head-of-line Blocking

- QUIC knows about (HTTP/3) streams
- Processes them independently
- Means there is no strict ordering between streams anymore!
- In HTTP/2: arrives in the order you put it on the wire, even if on different streams
- In HTTP/3: arrives ordered per-stream, but not across streams
 - E.g., if packet 1 for stream A, which was sent first, is lost, packet 2 for stream B, which was sent second, will arrive before retransmit of 1

►► UHASSELT EDN

4. HTTP/3 vs HTTP/2 : Exclusive priorities



https://tools.ietf.org/html/draft-ietf-quic-transport-19 https://tools.ietf.org/html/draft-ietf-quic-http-19 https://speeder.edm.uhasselt.be/www18/files/h2priorities_mwijnants_www2018.pdf



4. HTTP/3 vs HTTP/2 : Exclusive priorities

What if you add 2 nodes B and C exclusively to A at the same time?



"Solution" 1: No more exclusive priorities in HTTP/3 "Solution" 2: Send priority updates on 1 "control stream"

(only updates, not the initial, because HOL blocking!)

►► UHASSELT EDM

https://tools.ietf.org/html/draft-ietf-quic-transport-19 https://tools.ietf.org/html/draft-ietf-quic-http-19 https://speeder.edm.uhasselt.be/www18/files/h2priorities_mwijnants_www2018.pdf https://github.com/quicwg/base-drafts/issues/2502

4. HTTP/3 vs HTTP/2 : QPACK vs HPACK

- HTTP header compression
- HPACK: just send encoding information with the header
- QPACK:
 - Separate encoder and decoder streams
 - Either accept HOL-blocking, or keep sending literals until confirmed
 - Cannot just send with header, because other streams might start using encoded value and arrive before the encoding info...



4. HTTP/3 vs HTTP/2 : Priorities (again)



UHASSELT EDM

HTTP/2: "fake" streams as grouping nodes HTTP/3: explicit "placeholder" support built-in

> https://tools.ietf.org/html/draft-ietf-quic-transport-19 https://speeder.edm.uhasselt.be/www18/files/h2priorities_mwijnants_www2018.pdf https://tools.ietf.org/html/draft-ietf-quic-http-19

4. HTTP/3 vs HTTP/2 : Push

- PUSH_PROMISE
 - HTTP/2: STREAM_ID
 - HTTP/3: PUSH_ID (more flexible)
- DUPLICATE_PUSH (again: saves on bytes)
 - Due to reordering, DUPLICATE_PUSH frames can arrive before the corresponding PUSH_PROMISE frame

5. Additional stuff I could talk about

- Max_uni/bidi transport parameters + asymmetric streams
- Flow control
 - Connection vs stream-level
- Recovery
 - Combination of various best practices for loss detection
 - Pluggable congestion control
- Linkability
- Prioritization and buffer bloat
- QUIC standardized logging and debugging tools

